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Identification

Directory Supervisor, System Primitives
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Purpose

The directory supervisor provides the primitives for manipulating directory entries and decides the permission needed to carry out the requested operation from the intent of the caller.

Primitive

Certain primitives of the directory supervisor are privileged, i.e., they are callable only by procedures residing in system rings, namely the hard-core ring and the administrative ring. Some of these primitives are necessary to cleanup and update certain file system tables during activation, deactivation or reactivation of a segment (finddir, activinfo, refindb) and another to set up certain values for variables in system data bases (set_base_dir). Since these functions are necessary for the well-being of the file system, the question of whether or not the user has any attribute (read, execute, write, append) ON with respect to the segment being referenced to complete the function is irrelevant. In the terminology of directory control, the user needs a null access mode with respect to the affected segment.

In the case of estblseg and setusage, the access mode of the user is important. In order to establish a non-directory segment the user must have the execute attribute ON for the directory containing the segment's branch but he doesn't need any attribute ON for the segment itself. Since, for example, the user may be expecting a mode change with respect to this segment at any moment, he is allowed to establish a non-directory segment to which he has no access. In order to establish a directory segment the user may have a null access mode with respect to the directory containing this segment's branch but he needs at least one attribute ON with respect to the segment itself. This is used to restrict a user to a section of the file hierarchy beneath a particular node (directory).

In order to set the usage status of a particular segment the user needs the execute attribute ON in the directory

containing the segment's branch. He also needs that attribute ON with respect to the segment which corresponds to the usage being set; the read attribute if setting read usage, write or append attribute if setting write usage, or read, write or append attribute if setting data-share usage.

The following is a list of the primitives with the access mode needed in the directory containing the segment's branch and possibly the access mode needed in the branch (for the segment) itself.

1. `estblseg` (non-directory segment - execute
directory segment - at least one attribute
ON in branch)
2. `setusage` (execute; read in branch for read usage, write
or append in branch for write usage, or
read, write or append in branch for data
share usage)
3. `set_base_dir` (null)
4. `finddir` (null)
5. `refindb` (null)
6. `activinfo$rdbranch,activinfo$wrbranch` (null)

1. estblseg

The primitive estblseg is privileged and is called by the Segment Management Module after it has received an initiate call from a user. This primitive finds the branch to which a given entry points and, through a call to segment control at `makeknown` makes the segment known to the current process.

```
call estblseg (dir, entry, segsw, segptr, uid, optsw,
              slotlist, code);
```

```
dcl dir char(*), /*path name of a directory*/
```

```
entry char(*), /*name of entry in dir which points
                to the requested segment*/
```

```
segsw fixed bin(1), /*switch indicating whether the
                    user does (=1) or doesn't
                    (=0) want a special segment
                    number given to this segment*/
```

```

segptr ptr, /*pointer to the base of the established
            segment, returned by directory supervisor
            or given by caller if segsw is ON*/

uid bit(70), /*unique identifier of the segment
            found in the branch pointed to by
            entry*/

optsw bit(2), /*value of the option switches in the
            branch pointed to by entry, returned
            by directory supervisor*/

slotlist(*)fixed bin(17), /*list of slot numbers for
            each of the branches in
            the path name of the branch
            pointed to by entry,
            returned by directory
            supervisor*/

code fixed bin(17); /*if non-zero, it represents the
                    code of an error detected by the
                    file system*/

```

If the branch to which entry points is a non-directory branch then the user needs the execute permission in the directory containing this branch. If it is a directory branch, the user needs no permission in the directory which contains it.

The `findbranch` primitive is called to find the branch to which `entry` points. The `effmode` primitive is called to determine the effective mode of the user with respect to the segment. If the segment is a directory segment then the user must have at least one permission (read, execute, write or append). Segment control is called at `makeknown` to make the segment known and to return the segment pointer and slot list for the segment. The branch is then unlocked and control is returned to the caller.

2. setusage

The primitive setusage is privileged and is called by the Segment Housekeeping Module after it has received a segment usage request from a process. This primitive changes the usage of a given segment by this user from a previous usage to a new usage status if this new status is compatible with the usage of the segment by other processes.

```

call setusage (dir, entry, newuse, olduse, blocksw,
              blockrtn, code);

```

```

dcl dir char(*), /*path name of a directory*/
     entry char(*), /*name of the entry in dir which points
                    to the requested segment*/

newuse bit(2), /*designation of the new way in which
               the process wishes to use the segment;
               it may designate read (01), write
               (10), data-share (11) or no use (00)*/

olduse bit(2), /*designation of the current usage of
               the segment by this process*/

blocksw fixed bin(1), /*a switch indicating whether
                      a process is (ON) or isn't
                      (OFF) willing to wait until
                      the value of newuse becomes
                      compatible with the usage
                      of the segment by other
                      processes*/

blockrtn label; /*place to which control will be
                 returned if blocksw is ON and a
                 blocking situation occurs (see
                 below)*/

```

The user needs the permission implied in newuse. The primitive setusage first calls findbranch to find the branch to which entry points.

If olduse indicates that the process was using this segment for reading, writing or data-sharing, then the effect of the usage of this segment by this process must be removed from the usage and usage-count items in the branch for the segment. That is, the usage-count item is decreased by one and, if it is still greater than zero, the usage item is unchanged. Otherwise if the usage-count turned zero, then the usage item in the branch is changed to not-used and if the no-more users switch is ON, it is turned OFF and a call is made to notify in the process wait and notify (PWN) module to wake up any processes that are blocked waiting for this segment to become available.

If olduse indicates that the process was not using the segment then the usage and usage-count items are not changed with respect to this argument.

If newuse indicates that the process wishes to use the segment for reading, writing or data-sharing, then the usage and usage-count items in the branch must be updated to reflect this. If blocksw and the no-more-users switch are both ON, a call is made to addpwt in the PWN module to put this process on the process-waiting table for the correct event (so that it may later be notified if the event has occurred) and control is returned to blockrtn. If the no-more-users switch is OFF and

1. if the current usage item indicates not-used, then it is set to newuse and the usage-count item is set to 1, or
2. if the usage item equals newuse and is either read or data-share, then the usage-count item is increased by 1.

For any other combination, if blocksw is ON, the no-more-users switch is turned ON, addpwt is called and control is returned to blockrtn or, if blocksw is OFF, an error is reflected to the calling procedure.

3. set base dir

The primitive set base dir is privileged and is called by the file system initializer procedure initialize_branches (see BL.3) and certain administrative procedures. This primitive sets all file system activity for a process to a particular path from the root of the hierarchy.

```
call set_base_dir (pathname, code);
```

```
dcl pathname char(*); /*pathname of the particular
                        subtree to be used for file
                        system activity, given by
                        caller*/
```

No permission is needed to switch the subtree using set_base_dir since this is a privileged primitive.

The contents of pathname which for example may be

1. "root>system_root", the subtree used exclusively for segments of the trigger, hardcore supervisor and the hierarchy reconstruction process (see BH.3.01), or
2. "root>multics_root", the subtree used for all other segments,

are stored in the process definition segment for the current process. The size of pathname is also stored in the process definition segment. Any subsequent call to directory supervisor which specifies the path name of a directory relative to the root (the normal type of path name handed to directory control, see BX.8.00) has pathname prefixed to it. Hence all file system activity after this call is switched to the hierarchy subtree defined by pathname.

4. finddir

The primitive finddir is called only by segment control. It locates the directory branch to which a given entry effectively points and makes the segment to which this branch points known to the current process. This primitive is used by segment control when it is asked (by directory supervisor or maintainer) to return a segment pointer to a directory which is not known to the process.

```
call finddir (dir, segptr, code);
```

```
dcl segptr ptr; /*pointer to the segment with name dir,
                returned by directory supervisor*/
```

No permission is needed to find a directory and make it a known segment using this primitive. If dir does not denote the root directory, then findbranch is called to find and lock the branch which points to dir. If dir is not a directory, its branch is unlocked and an error is reflected to the caller.

If dir does denote the root directory (i.e., dir = "root") then the primitive findbranch need not be called because the branch which points to the root is known to the directory supervisor.

In either case, segment control is called at makeknown to make dir known to the process. The branch is unlocked and control is returned to the caller.

5. refindb

The primitive refindb is privileged and used only by segment control. It returns to segment control selected information from a requested branch. This primitive is used by segment control when activating and reactivating known segments for a process.

```

dcl dp ptr, /*pointer to the base of the directory segment
            containing the desired branch, given by caller*/

slot fixed bin(17), /*slot number of the branch when
                    originally found by estblseg, given
                    by the caller*/

uid bit(70), /*unique identifier of the branch when
             originally found by estblseg, given by
             the caller*/

dtbm bit(52), /*date and time this requested branch
              was last modified, returned*/

emode bit(5), /*effective mode of the current user with
              respect to the branch, returned*/

plist(plistmax)bit(18); /*list of protection numbers
                        defining access bracket, call
                        bracket and gates (see BG.9.00)
                        for the segment pointed to by
                        this branch (plistmax is a static
                        variable initialized to the
                        maximum size of any protection
                        list), returned*/

```

The user needs no permission to find a given branch using this primitive since it is privileged. The primitive `refindb` first checks the unique identification of the branch defined by `dp` and `slot` to see if it matches the given `uid`. If there is no match an error is reflected to the caller. If there is a match the `dtbm`, `emode` and `plist` items are taken from the branch and returned to the caller. The `effmode` primitive of access control must be called to compute the effective mode (`emode`) and the protection list (`plist`) from the access control list for the branch.

6. activinfo\$rdbranch, activinfo\$wrbranch

The primitives `rdbranch` and `wrbranch` are privileged and are used only by segment control. `rdbranch` is used by segment control to obtain the information from a given branch which is needed to activate a segment. `wrbranch` is called to update this information when the segment is deactivated or when segment control is asked specifically to update it.

```
call activinfo$rdbranch (dp, slot, uid, actsw, itemsptr, fmp
code);
```

```
call activinfo$wrbranch (dp, slot, uid, actsw, itemsptr, fmp
code);
```

```
dcl dp ptr, /*pointer to a directory segment, given by
caller*/
```

```
slot fixed bin(17), /*slot number of a branch in the
directory, given by caller*/
```

```
uid bit(70), /*unique identifier of the branch, given
by caller*/
```

```
actsw bit(1), /*new setting for the file-active switch
in the branch, given by caller*/
```

```
itemsptr ptr, /*pointer to a structure to be filled in
with certain items from the branch for
rdbranch or a structure containing new
values for these certain items for
wrbranch*/
```

```
fmp ptr; /*pointer to the file map of the branch*/
```

The user does not need any permission to read from or write into the branch using these primitives since they are privileged.

If rdbranch was called, then the requested branch must be locked. Before reading or writing the information which is pertinent when a segment is to be activated or deactivated, the secondary lock in the branch which locks only this subset of information must be set. When this lock is set, this information may be read or written. The file-active switch is set to actsw, all locks which were set are removed and control is returned to the caller.