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Identification

BCPL-MULTICS compiler - code generation
R. H. Canaday

Purpose

This document discusses the code produced by the BCPL-MULTICS compiler. Most of this document consists of descriptions of various detailed aspects of the code. The reader is presumed to be familiar with Ref(1) and with MSPM sections BZ.6.00 through BZ.6.03.

This document is not intended for the casual reader. It will be helpful to refer frequently to an EPLBSA listing produced by the BCPL compiler (such as is provided in the appendix) while reading the following discussions.

Summary

This document is subdivided as follows:

1) Initialization

Unlike EPL, BCPL programs cannot always be initialized on first reference. For example, a BCPL procedure cannot communicate through the global vector with procedures in other segments unless those segments have been initialized.

2) Stack management

Since BCPL does not use a standard MULTICS call-save-return for intercommunication, a separate BCPL stack is set up and maintained. The programmer has a great deal of (optional) control over the stack and global vector.

3) Call-Save-Return

Intercommunication between BCPL procedures is described.

4) Communication with EPL Procedures

Run-time routines have been provided which make possible recursive calls from EPL to BCPL procedures and vice-versa.

5) How to read EPLBSA produced by the BCPL compiler.

The listing from a compilation is not easy to correlate with the source code which produced it. However, the details given in this section should make the task somewhat easier.

Initialization

A BCPL procedure segment name is the same as the name of the CTSS or MULTICS source file which produced it. Unlike EPL, the segment name has nothing to do with the name of any procedure within the segment.

A BCPL procedure segment should be initialized before it is executed. The initialization of a segment named 'XXX' can be done either by calling (from EPL) the initialization entry point 'XXX\$XXX' [e.g., 'call XXX'] or by calling (from EPL) a procedure within the segment [e.g., 'call XXX\$PRO2 (arg1,arg2,...)'].

Example:

A program consists of three BCPL procedure segments named 'ASEG', 'BSEG', and 'CSEG'. The entry point is the procedure 'PRO3' in 'BSEG'. There are no arguments. The EPL driver would be as follows:

```
START:  PROC ( );
        CALL ASEG;
        CALL CSEG;
        CALL BSEG$PRO3;
        END;
```

In BCPL there are two forms of static storage which must be initialized: 'global' and 'local'. Local static storage is kept in the linkage segment and is truly static - like EPL static it is initialized only once in a given process. Repeated calls to initialize it will have no effect. A switch word, "INITSW", in each linkage segment is set to zero to indicate that initialization has been done. Global static, on the other hand, is kept in the 'global vector' (see the next section, Global and Stack Management, for details on the allocation of the global vector). Global static can be reinitialized repeatedly. Calling a BCPL procedure ('call BSEG\$PRO3') will initialize global static only if it has never been initialized for the segment BSEG (i.e., if INITSW \neq 0), but calling the initialization entry point ('call BSEG') will always reinitialize global static for BSEG. This is used to prevent interference between logically separate BCPL programs running in the same process.

Segment `BCPLGL` contains the routines which initialize static storage and which accept EPL calls to BCPL procedures and do argument conversion. The entry points are:

- `BCPLGL\$GINIT` called automatically by BCPL procedures to initialize global static and, if necessary, local static.
- `BCPLGL\$GSETU` called automatically by BCPL procedures to do stack setup and argument conversion.

There are also other entry points which are described in later sections of this document.

Global and Stack Management

Since BCPL does not conform to MULTICS standards in its stack management, it cannot use `SP` to point to its current stack frame. `SP` always points to a valid MULTICS stack frame. While BCPL procedures are in execution, `SP` points to a stack frame immediately inferior to (called by) that of the caller of BCPL, which contains only the header information and two words of diagnostic information: SP|32 contains an ITS pair which identifies the BCPL procedure segment which was called from EPL, and SP|34 contains an ITS pair which points to the global vector and the stack frame of the first BCPL procedure that was called.

The BCPL global vector and stack are contained in a segment which may be supplied by the caller or routines superior to it, or which will be created during BCPL initialization if not supplied. BCPL uses "unique_chars" and "smm\$set_name_status" to create a segment of 250K words. The global vector, which has a standard length of 2048 words, always starts at word 0 of the segment. The stack area begins immediately after the last word of the global vector.

During execution of a BCPL procedure the address register pairs are used as follows:

- AP - points to current BCPL stack frame
(AB points to global vector)
- BP - general usage
- SP - MULTICS stack frame
- LP - linkage segment

The registers are always paired. AB can be used for the global vector because the global vector always begins at word 0 of the segment containing the BCPL stack frames.

BCPL stack segment management centers around a pointer called the "Current Start of Stack" pointer, abbreviated "CSS". This pointer is kept in static storage and is accessed and changed through four routines. Routine BCPLGL\$SETAP can be called from EPL to set CSS to any address, including null. Routine BCPLGL\$LOOKAP returns in its argument the current value of CSS. Both entries have one argument, of type 'PTR'.

The third routine, BCPLGL\$GSETU, is not callable from EPL. It is invoked automatically whenever an EPL procedure calls a BCPL procedure. Its function is to do EPL-BCPL argument conversion and to load the address pair AB-AP with the value of CSS.

The fourth routine is the BCPL-callable procedure 'Call', accessed through global vector location 14. Its effect on CSS is to save the old value of CSS and reload CSS with a pointer to the first free word after the current BCPL stack frame. 'Call' then calls the desired EPL procedure. On return the current value of CSS is discarded and the saved value is restored.

The commonest use of these routines is to make sure that two independent BCPL programs (sets of procedures) running in the same process do not conflict in their use of the global vector. The simplest way to guarantee that a BCPL program will be executed with a clean global-and-stack segment is to issue the call

```
CALL BCPLGL$SETAP (NULL)
```

before initializing any BCPL procedure segment in the program. Note that any 'CALL BCPLGL\$SETAP' issued after the initialization of a BCPL procedure segment will invalidate that initialization. The only exception to this is that if a BCPL procedure calls an EPL procedure, nothing done by that procedure or any inferior procedure can affect that value of CSS which will be restored when the EPL procedure returns to the BCPL procedure. In fact, a good way of insuring that the global vector will be unchanged by (and invisible to) inferior procedures is to issue the call

```
CALL BCPLGL$SETAP (NULL);
```

in any EPL procedure called from a BCPL procedure.

In the case of a BCPL-to-EPL-to-BCPL calling chain in which the inferior BCPL procedure should reference the same global vector as the superior one, the superior routine can protect itself from changes made by the inferior routine by issuing the BCPL calls:

Save ()

before calling the EPL routine, and

Restore ()

on return. The effect of 'Save ()' and 'Restore ()' is to save and restore the global vector in a pushdown stack (which is kept in the MULTICS stack). The call 'Save ()' does not change the contents of the global vector. 'Save ()' and 'Restore ()' can be used not only to bracket a call to an EPL procedure, but in fact at any point in a BCPL procedure.

The routines 'BCPLGL\$SETSP(PTR)', 'BCPLGL\$LOOKAP(PTR)', 'Save ()', and 'Restore ()' give the programmer complete control over global static storage. However, he has very little control over local static storage. Thus in the MULTICS environment it may be useful to use portions of the global vector for static storage. However, be warned that these routines for controlling global static may be dependent on the MULTICS environment. Thus coding dependent on them may not be machine independent.*

Call-Save-Return

The BCPL stack frame header consists of 8 words, containing the following information:

<u>WORD</u>	<u>CONTENTS</u>	<u>STORED BY</u>
0-1	RETURN ITS	SAVE
2-3	LP FOR CURRENT PROCEDURE	SAVE
4	UNUSED	
5 UPPER	LEFT-HAND-SIDE FLAG	CALL
5 LOWER	ARGUMENT COUNT	CALL
6-7	AP FOR SUPERIOR PROCEDURE	SAVE

The code produced by the compiler for call-save-return in procedure "PROCED" is as follows. In this case the new stack frame will start at the nth word of the previous frame.

<u>CALL</u>		
LDA	m,DL	# arguments
STA	AP n+5	
EAX2	n	
LDA	CALLADDRESS	A BCPL ADDRESS
EABBB	0,AU	
TSBBP	BE 0,AP	

* The CTSS (and GECOS) environments act as if a 'BCPLGL\$SETAP(NULL)' call were issued at every return to command level (between activities).

SAVE

in the linkage segment

CALLADDRESS:	EAPLP	NULL	
	EAX7	STAT-*, IC	Set LP
	TRA	<PROCED> SV	The procedure Save Sequence

 in the procedure segment ('PROCED')

SV:	NULL		
	STPAP	AP 6,2	old AP
	STPBP	AP 0	return point
	EAPAP	AP 0,2	new AP
	STPLP	AP 2	new LP
	SZN	LP INITSW	zero if static has been initialized
	TZE	0,7	enter procedure body
INIT:	NULL		initialization procedure

RETURN

EAPAP	AP 6,*	old AP
EAPLP	AP 2,*	old LP
TRA	AP 0,*	return

Communication with EPL Procedures

As has previously been stated, communication between BCPL procedures is by means of a special calling sequence rather than the standard MULTICS call. Some of the significant differences between a BCPL call and a standard MULTICS call are:

- 1) BCPL arguments do not have dope
- 2) BCPL arguments are by value.
- 3) The BCPL stack frame header is 8 words long rather than 32.
- 4) Registers are not saved and restored across a call, except for AB-AP, LB-LP, and SB-SP.
- 5) SB-SP is not used. The stack pointer is AB-AP.
- 6) The address of the callee is stored in a variable.

The link-by-name features of MULTICS are not used. Obviously, despite these differences, a mechanism must exist to permit calls from EPL procedures to BCPL procedures and vice-versa.

In order to enable EPL procedures to call BCPL procedures the BCPL compiler generates two entry points for each BCPL procedure. One of these is used in BCPL-to-BCPL calls

to the procedure. The other entry point, labeled with the name of the procedure, is callable from EPL. This entry point invokes the routine `BCPLGL\$SETU` which transforms the call into BCPL format and generates a BCPL argument list from the EPL argument list. Dope and specifiers are not interpreted. The addresses in the EPL arglist are translated into BCPL addresses and placed in the BCPL arglist. Dope and specifiers can be interpreted by the callee as desired. Two functions exist to aid in interpreting arguments. These are `MtoBaddr(X)` which accepts the (BCPL) address of an ITS pair and returns the value of the ITS pair as a BCPL address, and `MtoBstring(X,V)` which accepts the (BCPL) address `X` of a MULTICS string specifier and places a corresponding BCPL string in vector `V`.

Calling from a BCPL procedure to an EPL procedure is done using the two library routines `Getadr` and `Call` (cf MSPM BZ.6.02). All of the arguments of `Call` are addresses. The first argument is the address to be called. The remaining addresses will be converted to ITS pairs and stored in an EPL-format argument list. If any arguments require dope and specifiers, they must be generated by the caller before calling `Call`. Two library functions, `ITS` and `BtoMstring`, are provided for this purpose. Function `ITS` creates an ITS pair from a BCPL address. Function `BtoMstring` creates a fixed length MULTICS character string from a BCPL string. MSPM BZ.6.02 contains more detail on `ITS` and `BtoMstring`.

The only complicated part of the method for calling between EPL and BCPL procedures is the stack management, which was described in detail above. In general it is possible to call freely between EPL procedures and BCPL procedures if one remembers that in the absence of calls to `BCPLGL\$SETAP`, there will not be any conflicts in stack usage. It is also worth noting that the EPL call

```
CALL BCPLGL$SETAP (NULL);
```

is always a safe way of guaranteeing that BCPL programs which precede and follow the call will be independent.

How to read EPLBSA produced by the BCPL compiler

The BCPL compiler produces some comments to help correlate the EPLBSA code with the source program: the source-program name of each variable is given at the point it is declared; the source name of each function and variable labels the corresponding entry point; and each call to a function or routine is labelled, at its occurrence in the code, with the source name of the callee.

Register usage by the code generators is as follows:

A	used for addresses or computations
Q	
AQ	used for floating point, multiply, and divide
x3-x4	always used as a pair - x4 = upper half word, x3 = lower
x5-x6	a pair like x3-x4
x2	contains the stack increment (frame size) during a call
x0	used with BB for indirection
x1,x7	unused

Multiple location counters are used, as follows:

MAINC	procedure segment main counter
LASTC	procedure segment builtin subroutines
STATC	linkage segment static storage to be initialized
LSTATC	linkage segment static storage not needing initialization
IGLØBC	procedure segment data to be put in the global vector

The next part of this discussion consists of examples of BCPL statements and the code generated by them, together with a brief description where necessary.

<u>Example A</u>	declarations	
let a,b,c = 0,0,5	STZ AP 6+10	"Declare a
	STZ AP 6+11	"Declare b
	LDA 5,DL	
	STA AP 6+12	"Declare c
let v = vec 100	EAPBP AP 6+14	
	EAX3 AP 6+14	
	SBRBB AP 6+13	
	SXL3 AP 6+13	"Declare v

This is an example of a very common operation, namely address generation. The variable 'v' is at location 19 of the stack frame. The vector begins at location 20. '19' is written '6+13' because of a historical carry-over from GMAP.

Example B

assignments

a :=v[20]

```
LDA AP|6+13    v
EABBB 0,AU
LDA BB|20,AL
STA AP|6+10    a
```

This is an example of using an address. Addresses may be used from registers A or Q, or from register pairs x3-x4 or x5-x6.

v[20] := b[a]

```
LDQ AP|6+13    v
LDA AP|6+11    b
ADA AP|6+10    a
EABBB 0,AU
LDA BB|0,AL
EABBB 0,QU
STA BB|0,QL
```

Example C

address arithmetic

v[0] := v[20] * a[b+c]

```
LDA AP|6+10    a
ADA AP|6+11    b
ADA AP|6+12    c
EAX4 0,AU
EAX3 0,AL
LDQ AP|6+13    v
EABBB 0,QU
LDQ BB|20,QL   v[20]
EABBB 0,4
MPY BB|0,3     a[b+c]
EABBB AP|6+13  v
STA BB|0,8
```

Example Da string (does not need initialization)

```
L5:  USE
      NULL
      ØCT
      ØCT
```

```
LSTATC
005141142143
144145000000
```

an address (this will be transformed into a BCPL address which will be stored in place of the EAPBP)

	USE	STATC
L6:	NULL	
	EAPBP	LP/L8 label in linkage seg.
L7:	NULL	
	EAPBP	AP/L9 label in procedure seg.

an address to go in global vector location 20

	USE	IGLØBC
	EAPBP	AB/L10 label in procedure seg.
	ZERO	0,20

APPENDIX

BCPL invoked through Mrgedt
Control Card = old bcpl test99

```
1 // This is a demonstration of BCPL code generation,  
2 // The program does not do anything and cannot be executed,  
3  
4 global $( Demonstration:100; Function:101 $)  
5 let Demonstration (A1,A2) = valof $(  
6     let a,b,c = 0  
7     let v = vec 100  
8     a := v[20]  
9     v[20] := b[a]  
10    v[0] := v[20] * a[b+c]  
11    c := "This is a string"  
12    a := Function(a, A1)  
13    Label:  
14    Label := Label + A2  
15    resultis A1 / Label $)
```

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OUTPUT FROM EPLBSA ASSEMBLY

.EPLBSA. PACKAGE 10 VERSION, 15 MAY 68.

.EPLBSA. BEGIN COMPILATION.

.EPLBSA.	ASSEMBLY OF FILE	\$TEST99\$,	SEGMENT NAME IS	TEST99
	000000		1	
	000000		2	NAME TEST99
	000000		3	FILE TEST99
S	000000	000032	4	LINK SAVER,<TEST99>~SV
			5	USE LASTC
	000114		6	LAST: NULL
			7	USE LSTATC
OA	000016	000121 0000 00	8	INITSW: ARG INIT
			9	USE STATC
			10	EIGHT
AA	000013	000000 000000	11	BSS STAT,4 "TEMP STORAGE
			12	USE IGLOBC
	000110		13	IGLOB: NULL
			14	USE MAINC
	000000		15	START: NULL
			16	JOIN /LINK/STATC,LSTATC/TEXT/IGLOBC,LASTC
	000000	000000	17	ENTRY TEST99
	000000		18	TEST99: NULL
AA	000000	6 00022 3521 20	19	SAVE
AA	000001	2 00020 6521 00		
AA	000002	2 00040 3521 00		
AA	000003	2 77762 2521 00		
AA	000004	2 77740 3321 00		
AA	000005	6 00032 2501 00		
OA	000006	000121 7070 00	20	TSX7 INIT
AA	000007	6 00020 1731 20	21	RETURN
AA	000010	6 00010 0731 00		
AA	000011	6 00024 6101 00		
OA	000012	000110 7100 00	22	TRA L2

0A	000013	000110	7100	00	23
	000014				24
	000014				25
	000014				26
	000014				27
	000014	000014			28
	000014				29
0A	000014	000127	7070	00	30
					31
	000017				32
LA	000017	777761	3700	04	33
0A	000020	000015	6270	00	34
4A	000021	4 00032	7101	20	35
					36
	000015				37
AA	000015	0 00012	4501	00	38
AA	000016	0 00013	4501	00	39
AA	000017	0 00014	4501	00	40
AA	000020	0 00016	3531	00	41
AA	000021	0 00016	6231	00	42
AA	000022	0 00015	5431	00	43
AA	000023	0 00015	4431	00	44
AA	000024	0 00015	2351	00	45
AA	000025	000000	3130	01	46
AA	000026	3 00024	2351	05	47
AA	000027	0 00012	7551	00	48
AA	000030	0 00015	2361	00	49
AA	000031	0 00013	2351	00	50
AA	000032	0 00012	0751	00	51
AA	000033	000000	3130	01	52
AA	000034	3 00000	2351	05	53
AA	000035	000000	3130	02	54
AA	000036	3 00024	7551	06	55
AA	000037	0 00015	2351	00	56
AA	000040	000000	6240	01	57
AA	000041	000000	6230	05	58
AA	000042	0 00012	2361	00	59
AA	000043	0 00013	0761	00	60
AA	000044	0 00014	0761	00	61

```

TRA          L3
"
"
" MAX STACK USED = 0
      SECRET TEST99, DEMONSTRATION
      ENTRY  DEMONSTRATION
DEMONSTRATION: NULL
      TSX7    BXSETU
      USE     LSTATC
LUA:         NULL
      EAPLP   STAT=-8, IC
      EAX7    L4A
      TRA     LP SAVED, *
      USE     MAINC
LUA:         NULL
      STZ     AP*6+4    *DECLARE A
      STZ     AP*6+5    *DECLARE B
      STZ     AP*6+6    *DECLARE C
      EAPBB   AP*6+8
      EAX3    AP*6+8
      SBRBB   AP*6+7
      SXL3    AP*6+7    *DECLARE V
      LDA     AP*6+7
      EABBB   O, AU
      LDA     BB*20, AL
      STA     AP*6+4
      LDQ     AP*6+7
      LDA     AP*6+5
      ADA     AP*6+4
      EABBB   O, AU
      LDA     BB*0, AL
      EABBB   O, QU
      STA     BB*20, QL
      LDA     AP*6+7
      EAX4    O, AU
      EAX3    O, AL
      LDQ     AP*6+4
      ADQ     AP*6+5
      ADQ     AP*6+6

```

AA 000045	000000	3130	02	62	EABBB	0,QU	
AA 000046	3 00000	2364	06	63	LDQ	BB*0,0L	
AA 000047	000000	3130	14	64	EABBB	0,4	
AA 000050	3 00024	4021	12	65	MPY	BB*20,3	
AA 000051	0 00015	2201	00	66	LDX0	AP*6+7	
AA 000052	000000	3130	10	67	EABBB	0,0	
AA 000053	0 00015	7201	00	68	LYLO	AP*6+7	
AA 000054	3 00000	7564	10	69	STQ	BB*0,0	
4A 000055	4 00022	3531	00	70	EAPBB	LP*16	
4A 000056	4 00022	6231	00	71	EAX3	LP*16	FORM ADDRESS
AA 000057	0 00014	5431	00	72	SBRBB	AP*6+6	
AA 000060	0 00014	4431	00	73	SXL3	AP*6+6	
AA 000061	0 00012	2351	00	74	LDA	AP*6+4	
AA 000062	0 00174	7551	00	75	STA	AP*6+118	
AA 000063	0 00010	2351	00	76	LDA	AP*6+2	
AA 000064	0 00175	7551	00	77	STA	AP*6+119	
AA 000065	000002	2350	07	78	LDA	2,DL	
AA 000066	0 00171	7551	00	79	STA	AP*6+115	
AA 000067	000164	6220	00	80	EAX2	6+110	
AA 000070	1 00145	2351	00	81	LDA	AB*101	
AA 000071	000000	3130	01	82	EABBB	0,AU	
AA 000072	0 00000	3571	00	83	STCD	AP*0	
AA 000073	3 00000	2721	05	84	TSBBP	BB*0,AL	"CALL FUNCTION *
AA 000074	1 00000	2351	00	85	LDA	AB*0	
AA 000075	0 00012	7551	00	86	STA	AP*6+4	
000076				87	NULL		
4A 000076	4 00014	2351	00	88	LDA	LP*15	
AA 000077	0 00011	0751	00	89	ADA	AP*6+3	
4A 000100	4 00014	7551	00	90	STA	LP*15	

* NOTE: CALL, SAVE, AND RETURN AS SHOWN HERE USE "STCD" BECAUSE OF A HARDWARE BUG IN THE "TSBBP" INSTRUCTION.

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OUTPUT FROM EPLBSA ASSEMBLY

AA	000101	0	00010	2361	00	91	LDQ	AP*6+2
4A	000102	4	00014	5061	00	92	DIV	LP*L5
AA	000103	1	00000	7561	00	93	STQ	AB*0
0A	000104		000105	7100	00	94	TRA	L8
						95	USE	LSTATC
	000022					96	NULL	
AA	000022		020124	150151		97	OCT	020124150151
AA	000023		163040	151163		98	OCT	163040151163
AA	000024		040141	040163		99	OCT	040141040163
AA	000025		164162	151156		100	OCT	164162151156
AA	000026		147000	000000		101	OCT	147000000000
						102	USE	STATC
	000014					103	NULL	
AA	000014	1	00076	3521	00	104	EAPBP	AB*L7
						105	USE	MAINC
	000105					106	NULL	
AA	000105	0	00006	3501	20	107	EAPAP	AP*6,*
AA	000106	0	00002	3701	20	108	EAPLP	AP*2,*
AA	000107	0	00000	6101	00	109	RTCD	AP*0
	000110					110	NULL	
	000110					111	"	
	000110					112	"	
	000110					113	"	MAX STACK USED = 128
						114	USE	IGLOBC
	000110					115	NULL	
4A	000110	4	00017	3521	00	116	EAPBP	LP*L4
AA	000111		000000	000144		117	ZERO	0,100
AA	000112		000000	000000		118	ZERO	
AA	000113		000000	000000		119	ZERO	
						120	USE	STATC
AA	000015		000000	3520	00	121	EAPBP 0	
						122	USE	LASTC
	000114					123	NULL	
AA	000114	0	00006	2501	12	124	STPAP	AP*6,2
AA	000115	0	00000	3501	12	125	EAPAP	AP*0,2
AA	000116	0	00002	6501	00	126	STPLP	AP*2
4A	000117	4	00016	2341	00	127	SZN	LP*INITSW
AA	000120		000000	6000	17	128	TZE	0,7
	000121					129	INIT:	NULL

```

4A 000121 4 00016 2354 00 130
4A 000122 4 00016 4504 00 131
AA 000123 6 00000 6504 00 132
0A 000124 000110 6260 00 133
2A 000125 000014 6250 00 134
4A 000126 4 00050 2724 20 135
000127 136
AA 000127 6 00022 3524 20 137
AA 000130 2 00020 6524 00
AA 000131 2 00040 3524 00
AA 000132 2 77762 2524 00
AA 000133 2 77740 3324 00
AA 000134 6 00032 2504 00
0A 000135 000114 3520 00 138
AA 000136 6 00000 6504 00 139
4A 000137 4 00052 7404 20 140
141

```

EXAMPTU:

```

LDA LP*INITSW
STZ LP*INITSW
STPLP SP*0
EAX6 IGLOB
EAX5 STAT+4
TSBBP <BCPLGL>*[GINIT]
MULL
SAVE
EAPBP SV
STPLP SP*0
TRA <BCPLGL>*[GSETU]
END

```

RETURN BY TRA BB*0,7

NO LITERALS

ENTRY POINTS AND SEGDEF NAMES

```

5A 000140 000006 000000
2A 000141 000044 000001
AA 000142 015 104 145 155 142 DEMONSTRATION
AA 000143 157 156 163 164
AA 000144 162 141 164 151
AA 000145 157 156 000 000
5A 000146 000012 000000
2A 000147 000036 000001
AA 000150 006 164 145 163 143 TEST99
AA 000151 164 071 071 000
5A 000152 000020 000000
6A 000153 000000 000002
AA 000154 014 163 171 155 144 SYMBOL\TABLE
AA 000155 142 157 154 137
AA 000156 164 141 142 154
AA 000157 145 000 000 000
5A 000160 000025 000000
6A 000161 000030 000002
AA 000162 010 162 145 154 145 REL\TEXT
AA 000163 137 164 145 170
AA 000164 164 000 000 000
5A 000165 000032 000000

```


AA 000167	010 162 145 154	146	REL\LINK
AA 000170	137 154 151 156		
AA 000171	153 000 000 000		
5A 000172	000037 000000		
AA 000174	012 162 145 154	147	REL\SYMBOL
AA 000175	137 163 171 155		
AA 000176	142 157 154 000		
AA 000177	000000 000000		

EXTERNAL NAMES

AA 000200	005 147 163 145	148	GSETU
AA 000201	164 165 000 000		
AA 000202	005 147 151 156	149	GINIT
AA 000203	151 164 000 000		
AA 000204	006 142 143 160	150	BCPLGL
AA 000205	154 147 154 000		

NO TRAP POINTER WORDS

TYPE-PAIR BLOCKS

AA 000206	000004 000000
55 000207	000044 000040